

A Real Symmetric Tridiagonal Matrix With a Given Characteristic Polynomial

Gerhard Schmeisser
*Mathematisches Institut
Universität Erlangen-Nürnberg
D-91054 Erlangen, Germany*

Submitted by Ludwig Elsner

ABSTRACT

Given a polynomial $u(x) = x^n + a_{n-1}x^{n-1} + \cdots + a_0$, $a_\nu \in \mathbf{R}$, $\nu = 0, 1, \dots, n-1$, having only real zeros, we construct a real symmetric tridiagonal matrix whose characteristic polynomial is equal to $(-1)^n u(x)$. This is a complete solution to a problem raised and partly solved by M. Fiedler.

1. INTRODUCTION AND STATEMENT OF RESULTS

A standard result in matrix theory states that the characteristic polynomial of a real symmetric matrix has all its roots real. At the International Colloquium on Applications of Mathematics in Hamburg (July 1990) Professor M. Fiedler presented a paper on a converse problem he has been thinking about for many years (see also [1]):

Construct a real symmetric matrix for which a given normed polynomial with only real roots will be the characteristic polynomial.

Of course, the construction should require only a finite number of numerical operations. He presented the following nice and simple solution [1].

THEOREM A. *Let $u(x)$ be a monic polynomial of degree $n \geq 1$; let b_1, \dots, b_n be distinct numbers such that $u(b_k) \neq 0$ for $k = 1, \dots, n$. Set*

$v(x) = \prod_{k=1}^n (x - b_k)$, and define the n -by- n matrix $A = (a_{ik})$ by

$$\begin{aligned} a_{ik} &= -\sigma d_i d_k & \text{if } i \neq k, \\ a_{kk} &= b_k - \sigma d_k^2, & i, k = 1, \dots, n, \end{aligned}$$

where σ is a fixed nonzero number and d_k is a root of

$$\sigma v'(b_k) d_k^2 - u(b_k) = 0.$$

Then $(-1)^n u(x)$ is the characteristic polynomial of the symmetric matrix A . If λ is an eigenvalue of A , then $(d_i/(\lambda - b_i))$ is a corresponding eigenvector.

If $u(x)$ is a polynomial with n real distinct roots, then choosing b_i as any real numbers interlacing the roots of u , one can choose σ as $+1$ or -1 in such a way that d_i are all real and A thus real symmetric.

Fiedler also showed that by a modification of the construction the matrix A takes the form

$$A = \begin{pmatrix} b_1 & & & c_1 \\ & \ddots & & \vdots \\ & & b_{n-1} & c_{n-1} \\ c_1 & \cdots & c_{n-1} & d \end{pmatrix},$$

where nonzero off-diagonal entries may appear only in the last row and column.

An interesting aspect of these results is that they allow a lot of flexibility concerning the choice of the numbers b_k which determine the elements of the matrix A . On the other hand, to guarantee that the elements of A are all real, the choice of appropriate b_k 's may require some preliminary calculations like construction of a Sturm sequence and performing bisections. Furthermore, the case of a polynomial with real but multiple zeros seems not to be settled.

The purpose of this note is to present an alternative solution to the above problem of Fiedler which provides a real symmetric tridiagonal matrix for every polynomial with only real zeros (not necessarily distinct). It is based on a modification of the Euclidean algorithm. There may be still other interesting solutions, and the reader is encouraged to look for them.

Let us start with a

NOTATION. For a polynomial

$$f(x) = a_k x^k + a_{k-1} x^{k-1} + \cdots + a_0$$

of degree $k \geq 0$ we write $c(f) := a_k$ so that $f/c(f)$ is a monic polynomial having the same zeros as f .

MODIFIED EUCLIDEAN ALGORITHM. For

$$u(x) = x^n + a_{n-1}x^{n-1} + \cdots + a_0, \quad a_\nu \in \mathbf{R}, \quad \nu = 0, 1, \dots, n-1, \quad (1)$$

define

$$f_1(x) := u(x), \quad f_2(x) := \frac{1}{n}u'(x)$$

and proceed recurrently as follows: If $f_{\nu+1}(x) \not\equiv 1$, then dividing f_ν by $f_{\nu+1}$ with remainder $-r_\nu$, we obtain

$$f_\nu(x) = q_\nu(x)f_{\nu+1}(x) - r_\nu(x) \quad (2)$$

and define

$$(i) \quad c_\nu := c(r_\nu), \quad f_{\nu+2}(x) := \frac{r_\nu(x)}{c_\nu} \quad \text{if } r_\nu(x) \not\equiv 0,$$

$$(ii) \quad c_\nu := 0, \quad f_{\nu+2}(x) := \frac{f'_{\nu+1}(x)}{c(f'_{\nu+1})} \quad \text{if } r_\nu(x) \equiv 0.$$

If $f_{\nu+1}(x) \equiv 1$, we terminate the algorithm, defining $q_\nu(x) := f_\nu(x)$.

THEOREM 1. *The polynomial (1) has only real zeros if and only if the modified Euclidean algorithm yields $n - 1$ nonnegative numbers c_1, \dots, c_{n-1} , and in this case*

$$u(x) = (-1)^n \det(T - xI),$$

Proof. The form of q needs no explanation. Clearly, r is of degree at most $k - 1$. If

$$\xi_1 > \xi_2 > \cdots > \xi_k \tag{5}$$

are the zeros of g , then by (4) and the separation property of the zeros of the monic polynomials f and g we have

$$(-1)^\nu = \operatorname{sgn} f(\xi_\nu) = -\operatorname{sgn} r(\xi_\nu), \quad \nu = 1, 2, \dots, k.$$

Hence r must be of the degree $k - 1$, and have precisely one simple zero in each of the intervals $]\xi_{\nu+1}, \xi_\nu[$, where $\nu = 1, 2, \dots, k - 1$. Furthermore,

$$\operatorname{sgn} r(x) = \operatorname{sgn} r(\xi_1) = 1 \quad \text{for } x \geq \xi_1,$$

which shows that $c(r) > 0$. ■

As a converse statement we shall need

LEMMA 2. *Let g and h be monic polynomials of degrees k and $k - 1$, respectively, whose zeros are all real, and distinct and separate each other. For $c > 0$ and $a \in \mathbf{R}$ define*

$$f(x) := (x - a)g(x) - ch(x). \tag{6}$$

Then f is a monic polynomial of degree $k + 1$ with distinct real zeros which are separated by those of g .

Proof. Obviously f is a monic polynomial of degree $k + 1$, and $f(x)$ is positive for large x . Denoting and ordering the zeros of g as in (5), we must have $h(\xi_1) > 0$, since h is monic and has no zero to the right of ξ_1 . By (6) and the separation property of the zeros of g and h ,

$$(-1)^{\nu+1} = \operatorname{sgn} h(\xi_\nu) = -\operatorname{sgn} f(\xi_\nu), \quad \nu = 1, 2, \dots, k.$$

Hence f has an odd number of zeros in $]\xi_1, +\infty[$ and in each of the intervals $]\xi_{\nu+1}, \xi_\nu[$, where $\nu = 1, 2, \dots, k - 1$. But this is possible only if f has precisely one simple zero in each of these intervals and an additional zero in $]-\infty, \xi_k[$. ■

3. PROOFS OF THE THEOREMS

We start with the

Proof of Theorem 2. If the polynomial (1) has n distinct zeros, they are separated by those of u' . Now repeated application of Lemma 1 shows that the modified Euclidean algorithm runs through $n - 1$ steps of type (i), yielding $n - 1$ positive numbers c_1, \dots, c_{n-1} .

Conversely, if the modified Euclidean algorithm yields $n - 1$ positive numbers c_1, \dots, c_{n-1} , it can only run through steps of type (i), providing $n - 1$ monic polynomials $q_\nu(x)$ of degree 1 such that

$$f_\nu(x) = q_\nu(x)f_{\nu+1}(x) - c_\nu f_{\nu+2}(x), \quad \nu = 1, 2, \dots, n - 1, \quad (7)$$

where $f_n(x) =: q_n(x)$ is of degree 1 and $f_{n+1}(x) \equiv 1$.

Fixing $q_\nu(x)$ and c_ν ($\nu = 1, 2, \dots, n - 1$) and starting with $f_{n+1}(x)$ and $f_n(x)$, we may successively calculate $f_{n-1}(x), f_{n-2}(x), \dots$ from (7). Applying Lemma 2 in each step, we find that $f_1(x) = u(x)$ has n distinct zeros. ■

Proof of Theorem 1. Suppose u has only real zeros, and let ϕ be the greatest common divisor of u and u'/n . Clearly, ϕ is a monic polynomial, which we may assume to be of positive degree $n - k$, say, since otherwise we are in the situation of Theorem 2. Consequently,

$$u(x) = f_1(x) = \phi(x)\varphi_1(x) \quad \text{and} \quad \frac{1}{n}u'(x) = f_2(x) = \phi(x)\varphi_2(x),$$

where φ_1 and φ_2 are monic polynomials of degrees k and $k - 1$, respectively, whose zeros are real and distinct and separate each other. Furthermore, the modified Euclidean algorithm shows that for $\nu = 1, 2, \dots, k - 1$ the polynomial ϕ factors out in the recurrence relation (7) so that

$$f_\nu(x) =: \phi(x)\varphi_\nu(x), \quad \nu = 1, 2, \dots, k, \quad f_{k+1}(x) = \phi(x), \quad \varphi_{k+1}(x) \equiv 1$$

and

$$\varphi_\nu(x) = q_\nu(x)\varphi_{\nu+1}(x) - c_\nu\varphi_{\nu+2}(x), \quad \nu = 1, 2, \dots, k - 1. \quad (8)$$

Since (8) is equivalent to the division of φ_ν by $\varphi_{\nu+1}$ with remainder $-c_\nu\varphi_{\nu+2}$, we may apply Lemma 1 successively for $\nu = 1, \dots, k - 1$ and find

that the numbers c_ν are all positive and the degree of $f_{\nu+1}$ decreases by 1 if ν increases by 1. The calculation of f_{k+2} leads to a step of type (ii) which yields

$$c_k = 0, \quad f_{k+2}(x) = \frac{1}{n-k} \phi'(x).$$

Continuing with f_{k+1} and f_{k+2} , we are in the same situation as at the beginning and may therefore repeat the above argumentation. This way we obtain precisely $n - 1$ numbers c_ν , which are all nonnegative.

If the modified Euclidean algorithm yields $n - 1$ nonnegative numbers c_ν , the relations (7) must hold with monic polynomials q_ν of degree 1. Since $q_n := f_n$, the matrix T in (3) exists and has real elements. Now define $T_1 := T$, and let T_k for $k > 1$ be the submatrix obtained from T by canceling the first $k - 1$ rows and columns. For $\nu = 1, \dots, n$ denote by g_ν the characteristic polynomial of T_ν . Obviously,

$$\begin{aligned} g_n(x) &= -q_n(0) - x = -q_n(x) = -f_n(x), \\ g_{n-1}(x) &= q_{n-1}(x)q_n(x)c_{n-1} \\ &= q_{n-1}(x)f_n(x) - c_{n-1}f_{n+1}(x) = f_{n-1}(x). \end{aligned} \tag{9}$$

Expanding $\det(T_\nu - xI)$, where I is the identity matrix of the same size as T_ν , with respect to the first row, we find

$$\begin{aligned} g_\nu(x) &= [-q_\nu(0) - x]g_{\nu+1}(x) - c_\nu g_{\nu+2}(x) \\ &= -q_\nu(x)g_{\nu+1}(x) - c_\nu g_{\nu+2}(x). \end{aligned} \tag{10}$$

Comparing (9) and (10) with (7), we conclude that $f_\nu(x) = (-1)^{n-\nu+1}g_\nu(x)$ and hence

$$u(x) = f_1(x) = (-1)^n \det(T - xI).$$

Since T is real and symmetric, the zeros of u are all real. ■

4. EXAMPLES

For the polynomial $u(x) = x^5 - 5x^3 + 4x$ whose zeros are $0, \pm 1,$ and ± 2 the matrix T becomes

$$T = \begin{pmatrix} 0 & \sqrt{2} & 0 & 0 & 0 \\ \sqrt{2} & 0 & \sqrt{\frac{7}{5}} & 0 & 0 \\ 0 & \sqrt{\frac{7}{5}} & 0 & \sqrt{\frac{36}{35}} & 0 \\ 0 & 0 & \sqrt{\frac{36}{35}} & 0 & \sqrt{\frac{4}{7}} \\ 0 & 0 & 0 & \sqrt{\frac{4}{7}} & 0 \end{pmatrix}.$$

For the polynomial $u(x) = x^5 - x^3$ the modified Euclidean algorithm runs through two steps of type (ii) and we obtain

$$T = \begin{pmatrix} 0 & \sqrt{\frac{2}{5}} & 0 & 0 & 0 \\ \sqrt{\frac{2}{5}} & 0 & \sqrt{\frac{3}{5}} & 0 & 0 \\ 0 & \sqrt{\frac{3}{5}} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}.$$

REFERENCES

1. M. Fiedler, Expressing a polynomial as characteristic polynomial of a symmetric matrix, *Linear Algebra Appl.* 141:265–270 (1990).
2. E. Isaacson and H. B. Keller, *Analysis of Numerical Methods*, Wiley, New York, 1966.
3. M. Marden, *Geometry of Polynomials*, Amer. Math. Soc., Providence, 1966.

Received 11 October 1990; final manuscript accepted 18 July 1991