Extending Simple Tabular Reduction with Short Supports

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Constraints, GAC

- Suppose we have finite-domain variables x₁, x₂, x₃ with domains x₁:{1,..,11}, x₂, x₃:{1,..,10}
- Constraint: $(x_1 = x_2 \text{ OR } x_1 = x_3)$
- Generalised Arc-Consistency (GAC) requires that each value of each variable is contained in a satisfying tuple of the constraint
- To *establish* GAC: *x*₁ ≠ 11

Support

- Suppose we have finite-domain variables x₁, x₂, x₃ with domains x₁:{1,..,11}, x₂, x₃:{1,..,10}
- Constraint: $(x_1 = x_2 \text{ OR } x_1 = x_3)$
- Traditional definition of GAC *support*: a satisfying tuple of the constraint
- Value $x_1 \rightarrow 11$ has no support, and is deleted
- Value x₁→1 is *not* deleted because it has support (1, 1, 3) (for example).

Short Support

- The key idea used in this paper:
- Suppose a constraint can be satisfied by an assignment to a small subset of its variables
 - This assignment is a *short support*
- Exploit these short supports to maintain GAC more efficiently

Short Support – Example

- Consider the running example again
- Domains **x**₁:{1,..,11}, **x**₂, **x**₃:{1,..,10}
- Constraint: $(x_1 = x_2 \text{ OR } x_1 = x_3)$
- Short support: ($x_1 \rightarrow 1, x_2 \rightarrow 1$)
- Any *extension* of this short support to cover X₃ is a full-length support

- Assuming we always use values in the domain

• Supports $x_1 \rightarrow 1$, $x_2 \rightarrow 1$, and all values of x_3

Short Support – Explicit and Implicit

- Consider the running example again
- Domains **x**₁:{1,..,11}, **x**₂, **x**₃:{1,..,10}
- Constraint: $(x_1 = x_2 \text{ OR } x_1 = x_3)$
- Short support: ($x_1 \rightarrow 1, x_2 \rightarrow 1$)
 - *Explicitly* supports $x_1 \rightarrow 1, x_2 \rightarrow 1$

– *Implicitly* supports all values of x_3

Short Support

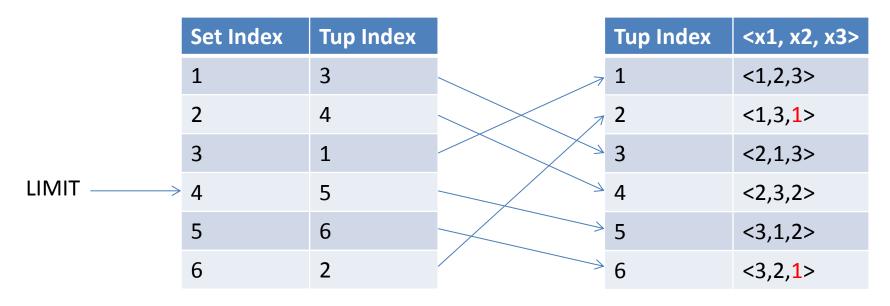


- Previously applied in GAC-Schema-like algorithms:
 - SHORTGAC (IJCAI 2011), then refined to HAGGISGAC (JAIR 2013)
 - HAGGISGAC is orders of magnitude faster than GAC-Schema when using short supports
 - HAGGISGAC a little faster than GAC-Schema with full-length supports (for an unrelated reason)
- Bigger goal: match the speed of hand-written propagators

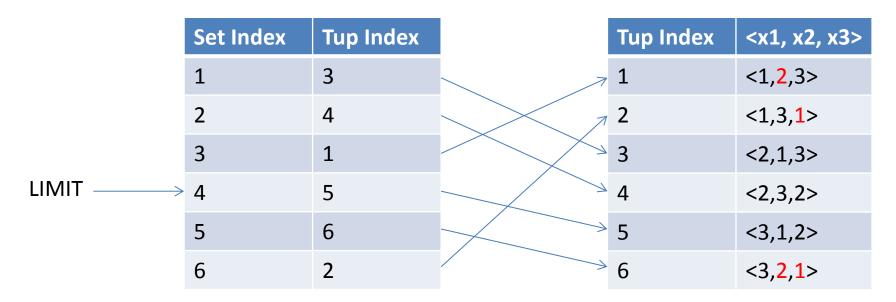
SHORTSTR2

- A new GAC algorithm extending STR2+ with short supports
 - Short supports are a perfect fit for STR2(+)
 - STR2(+) already optimises *fully supported* variables
 - The variable is removed from loops
- For each short support:
 - Variables with implicit support are marked as fully supported
 - Variable-value pairs with explicit support are treated exactly as in STR2+
- Given full-length supports, virtually identical to STR2+

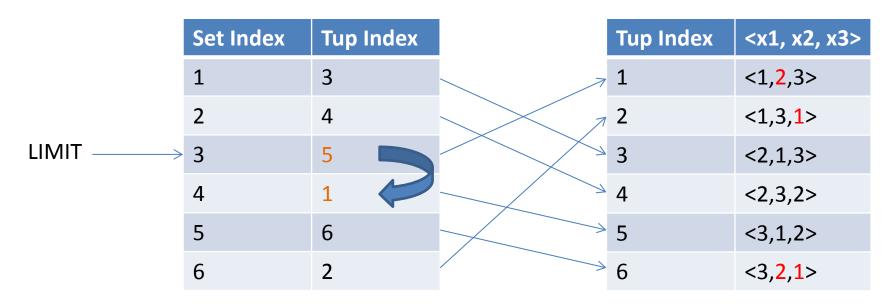
STR maintains a sparse set of the satisfying tuples



- Suppose x3, 1 is pruned
- Tuples 3,4,1,5 are in the set and 6,2 are out



- Now suppose x2, 2 is pruned
- STR algorithms iterate through tuples 3, 4, ...



- Now suppose x2, 2 is pruned
- STR algorithms iterate through tuples 3, 4, 1, ..
- Set now contains 3, 4, 5

- STR(2)(+) worst case complexity is terrible O(n²dⁿ⁺¹)
- Why are STR algorithms fast for some constraints?
- After just a few calls, set has been reduced enormously
- An extremely eager incremental propagator

Tuple Compression

- Take a set of full-length tuples and create a (non-unique) set of short supports
 – NP-hard to find a minimal set
- We propose a simple, fast greedy algorithm

Tuple Compression

- Using * to represent *any-value*
- Arity 4 constraint, each domain {1,2,3}
- Basic step is to take d (short) tuples and compress to one short tuple:

• Apply this rule to exhaustion

ShortSTR2 vs STR2+

- ShortSTR2 with tuple compression as a dropin replacement for STR2+
- Whole solver speed-up- ranges from 0.99 to 1.75

Problem class	Compression ratio	Speed-up ShortSTR2 compared to STR2+
Half	1.87	1.75
modifiedRenault	5.35	0.99
Rand-8-20-5	1.01	1.05
bddSmall	1.90	1.13
Renault	6.31	1.06
bddLarge	1.80	1.21
cril	1.19	1.11

Short Supports vs Full Length

- On Conway's Life and similar
- Problems are almost entirely one table constraint repeated
- Benefit of short supports varies

Problem	ShortSTR2 node rate Greedy compression	ShortSTR2 node rate Full length supports
Life	4,970	3,960
Brian's Brain	532	75
Immigration	4,930	3,590
QuadLife	483	>4GiB Memory

ShortSTR2 vs HaggisGAC

- Pigeonhole problem generalised to vectors of variables
- Vector not-equal constraints
- *p* is number of 'pigeons', *a* is number of variables per vector

р	а	ShortSTR2	HAGGISGAC
30	5	92,500	44,100
30	10	142,000	70,700
30	20	111,000	67,000
30	50	87,200	55,000
30	100	67,600	45,200
30	200	53,700	46,100

ShortSTR2 vs HAGGISGAC

- Pigeonhole problem generalised to vectors of variables
- Vector not-equal constraints
- *p* is number of 'pigeons', *a* is number of variables per vector
- Neither dominates the other complementary

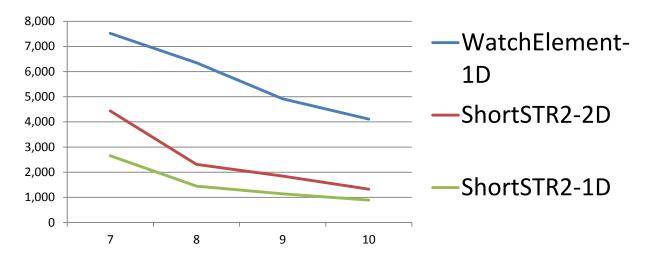
р	а	ShortSTR2	HAGGISGAC
5	100	592,000	1,790,000
10	100	250,000	653,600
20	100	119,000	158,800
30	100	67,600	45,200
40	100	43,700	18,000
50	100	31,900	10,900

ShortSTR2 vs HaggisGAC

- HAGGISGAC is orders-of-magnitude faster than Constructive Or and GAC-Schema (JAIR 2013)
 - When constraint is amenable to short supports
 - Element, Lex ordering, Square packing
- HaggisGAC approaches specialised propagators – particularly lex ordering

ShortSTR2 vs specialised propagator

• We compared to the Watched Element propagator on quasigroup problems



• 2x to 4x slower than hand-written propagator

Conclusions

- ShortSTR2 is a new GAC algorithm that extends STR2+ using short supports
 - Could be used as a drop-in replacement for STR2(+)
- Complementary to HAGGISGAC in performance

 Much simpler than HAGGISGAC
- Generic propagators as fast as specific handwritten ones?
 - Getting closer