

# Almost and Absolute Pure Acts over Semilattices

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## S-acts

Throughout,  $S$  is a monoid.

A right  $S$ -act is a nonempty set  $A$  together with a map

$$A \times S \rightarrow A, (a, s) \mapsto as$$

such that for all  $a \in A, s, t \in S$

$$(as)t = a(st) \text{ and } a1 = a.$$

- ❖ For any  $s \in S$ , we have an operation  $\rho_s : A \rightarrow A$  given by  $(a)\rho_s = as$ . The function  $\rho : S \rightarrow T_A$  given by  $(s)\rho = \rho_s$  is a monoid morphism.

Conversely, if  $\theta : S \rightarrow T_A$  is monoid morphism, define

$$as = (a)\theta_s$$

Then  $A$  is an  $S$ -act.

## Examples of S-acts

1.  $S$  is an  $S$ -act
2. Any right ideal of  $S$  is an  $S$ -act.
3. Let  $(K, +, \cdot)$  be a field and  $V$  be a left vector space. Then  $V$  is a left  $(K, \cdot)$ -act but not a  $(K, +)$ -act.
4. For any monoid  $S$  and a non-empty set  $A$ , define  $as = a$  for all  $a \in A$ , then  $A$  becomes a right  $S$ -act.

## Subact of S-acts

Let  $A_S$  be an  $S$ -act and  $B \subseteq A$ , a nonempty subst. Then  $B$  is a subact of  $A$  if  $as \in B$  for all  $a \in B$  and  $s \in S$ .

Obviously, any right ideal of  $S$  is a subact of  $S_S$ .

## Congruences and Morphisms for S-acts

- ❖ Let  $A$  be an  $S$ -act. An equivalence relation  $\sigma$  on  $A$  is called an  $S$ -act congruence or a congruence on  $A$ , if  $a\sigma b$  implies  $(as)\sigma(bs)$  for  $a, b \in A$  and  $s \in S$ .
- ❖ If  $X \subseteq A \times A$ , then  $\sigma(X)$  denote the smallest congruence on  $A$  containing  $X$ .
- ❖ A congruence  $\sigma$  is finitely generated if there exists a finite subset  $X \subseteq A \times A$  such that  $\sigma = \sigma(X)$ .
- ❖ The ordered pair  $(a, b) \in \sigma(X)$  if and only if either

$$a = b$$

or there exists a natural number  $n$  and a sequence

$$\begin{array}{ccccccc} a = c_1t_1 & & d_2t_2 = c_3t_3 & \cdots & & d_nt_n = b \\ d_1t_1 = c_2t_2 & & d_3t_3 = c_4t_4 & \cdots & & d_{n-1}t_{n-1} = c_nt_n \end{array}$$

where  $t_1, t_2, t_3, \dots, t_n \in S$  and for each  $i = 1, \dots, n$  either  $(c_i, d_i)$  or  $(d_i, c_i)$  is in  $X$ .

- ❖ If  $\sigma$  is a congruence on  $A$ , then  $A/\sigma$  is an  $S$ -act.
- ❖ An  $S$ -morphism from  $A$  to  $B$  is a map  $f : A \rightarrow B$  with  $(as)f = (af)s$  for all  $a \in A$  and  $s \in S$ .

## Free S-acts

- ❖ An S-act  $A$  is finitely generated if there exists a subset  $U$  of  $A$  such that

$$A = \bigcup_{u \in U} uS \text{ and } |U| < \infty.$$

- ❖ An S-act is free if there exists a subset  $U$  of  $A$  such that  $A = \bigcup_{u \in U} uS$  and each element  $a \in A$  can be uniquely presented in the form  $a = us$ ,  $u \in U$  and  $s \in S$ .
- ❖ Let  $X$  be a nonempty set. Then free S-act  $F(X)$  on  $X$  exists.

### Construction for $F(X)$ : Let

$$F(X) = X \times S$$

and define

$$(x, s)t = (x, st).$$

Then it is easy to check that  $F(X)$  is an S-act. With  $x \mapsto (x, 1)$ , we have  $F(X)$  is free on  $X$ .

Notice that

$$(x, s) = (x, 1)s \equiv xs.$$

- ❖ Free S-acts are disjoint unions of copies of  $S$ .

## Finitely Presented S-acts

An S-act  $A_s$  is cyclic is  $A = \langle \{a\} \rangle$ , where  $a \in A$ .

An S-act  $A$  is finitely presented if

$$A \cong F/\sigma$$

for some finitely generated free S-act  $F$  and finitely generated congruence  $\sigma$ .

Proposition:

Let  $A_s$  be a cyclic S-act. Then  $A_s$  is finitely presented if and only if it is isomorphic to a factor act of  $S_s$  by a finitely generated right congruence on  $S$ , that is,

$$A \cong S/\sigma$$

where  $\sigma$  is finitely generated right congruence.

## System of Equations over S-act

- ❖ Let  $A$  be an  $S$ -act. An equation over  $A$  has one of the three forms

$$xs = a \quad xs = xt \quad xs = yt$$

where  $s, t \in S$ ,  $a \in A$  and  $x, y$  are variables

- ❖ Let  $\Sigma$  be a system of equations over  $A$ . Then  $\Sigma$  is consistent if  $\Sigma$  has solution in some  $S$ -act  $B \supseteq A$ .

### Consistency criteria for $\Sigma$ :

Let  $\Sigma = \{ xs_i = a_i, xu_i = xv_j : s_i, u_j, v_j \in S, a_i \in A, 1 \leq i \leq n, 1 \leq j \leq m \}$  and  $\sigma = \langle (u_j, v_j) : 1 \leq j \leq m \rangle$ . Then  $\Sigma$  is consistent if and only if for all  $h, k \in S$  and for all  $1 \leq i, i' \leq n$ ,

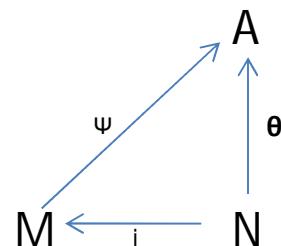
$$s_i h \sigma s_{i'} k \text{ implies } a_i h = a_{i'} k.$$

## Almost Pure S-act

An S-system A is almost pure if every finite consistent system of equations in one variable, with constants from A, has a solution in A.

Proposition: The following conditions are equivalent for an S-act A:

1. A is almost pure;
2. Given any diagram of S-acts and S-homomorphisms



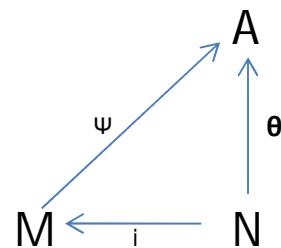
where M is cyclic finitely presented, N is finitely generated and  $i : N \rightarrow M$  is an injection, there exists an S-homomorphism  $\Psi : M \rightarrow A$  such that  $i \Psi = \theta$ ; further, for any  $s_1, \dots, s_n$  in S there is an element a in A with  $a = as_1 = \dots = as_n$ .

## Absolutely Pure S-acts

An S-system A is absolutely pure if every finite consistent system of equations, with constants from A, has a solution in A.

Proposition: The following conditions are equivalent for an S-act A:

1. A is absolutely pure;
2. Given any diagram of S-acts and S-homomorphisms



where M is finitely presented, N is finitely generated and  $i : N \rightarrow M$  is an injection, there exists an S-homomorphism  $\Psi : M \rightarrow A$  such that  $i \Psi = \theta$ .

## Completely Right Pure Monoids

A monoid  $S$  is called completely right pure if all right  $S$ -acts are absolutely pure.

Theorem: A monoid  $S$  is completely right pure if and only if all  $S$ -acts are almost pure.

Absolutely pure  $S$ -act  $\Rightarrow$  Almost Pure  $S$ -acts

For completely right pure monoids:

Almost Pure  $S$ -acts  $\Rightarrow$  Absolutely pure  $S$ -act

Does there exists an almost pure S-act which is not absolutely pure???

OR

Does there exists a class of monoids (Other than completely right pure) for which almost pure s-acts are absolutely pure????

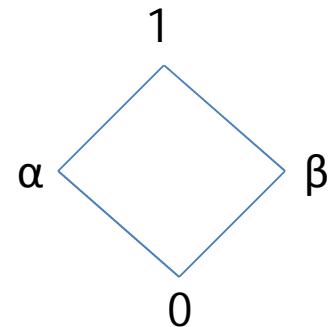
Theorem: A monoid  $S$  is completely right pure if and only if  $S$  has local left zeros and satisfies (\*):

(\*) given any finitely generated right congruence  $\sigma$  on  $S$  and any finitely generated right ideal  $I$  of  $S$ , there is an element  $s$  of  $I$  such that for any  $u, v$  in  $S$ , if  $u \sigma v$  then  $su \sigma sv$  and for any  $w \in I$ ,  $w \sigma sw$ .

Corollary: Let  $I$  be finitely generated right ideal of a completely right pure monoid. Then  $I = eS$  for some idempotent element  $e$  of  $S$ .

- ❖ The following conditions are equivalent for a monoid  $S$ :
  - i.  $S$  is regular and its principle right ideals are linearly ordered with respect to inclusion.
  - ii. Every finitely generated right ideal is generated by an idempotent element.

Consider the semilattice  $S=\{0, \alpha, \beta, 1\}$ .



- ❖  $S$  is an  $S$ -act over itself.
- ❖  $S$  is not completely right pure because the principle right ideals  $\{0, \alpha\}$  and  $\{0, \beta\}$  are not linearly ordered with respect to inclusion.
- ❖  $S$  is an almost pure  $S$ -act over itself.

For  $S=\{0, \alpha, \beta, 1\}$ , every almost pure  $S$ -act is absolutely pure.

Proof: Let  $A$  be an arbitrary almost pure  $S$ -act.

$S(n)$ : All finite consistent system of equations over  $A$  in no more than  $n$  variables have a solution in  $A$ .

Clearly, the assumption is true for  $n = 1$ .

Let  $\Sigma$  be a finite consistent system of equations in  $n+1$  variables  $x_1, x_2, \dots, x_n, x_{n+1}$ . Since  $\Sigma$  is consistent,  $\Sigma$  has a solution  $(b_1, b_2, \dots, b_n, b_{n+1})$  in some  $S$ -act  $B \supseteq A$ .

Case I:

$\Sigma$  has no equation of the form  $x_i = t$  for all  $i = 1, 2, \dots, n+1$ . Since  $A$  is almost pure  $S$ -act, so Solution of  $\Sigma$  exists in  $A$ .

Case II:

$\Sigma$  has an equation of the form  $x_i = t$  for some  $i = 1, 2, \dots, n+1$ . Since  $\Sigma$  is consistent. So  $b_i = a \in A$ .

Construct

$\Sigma' = \text{Considering all equations of } \Sigma \text{ in } x_2, \dots, x_n, x_{n+1} \text{ and}$   
 $\text{replacing } x_1 \text{ by } a.$

$\Sigma'$  is consistent. So, by induction  $\Sigma'$  has a solution  $(a_2, a_3, \dots, a_n, a_{n+1})$  in A.

**$(a_1 = a, a_2, \dots, a_n, a_{n+1})$  is the solution of  $\Sigma$  in A.**

### Case III:

Suppose  $\Sigma$  has equations of the form :

$x_{iS} = a$  and  $x_i = x_{i'}t$  for some  $i, i' = 1, 2, \dots, n+1$  and  $i \neq i'$

Take  $i = 1$  and  $i' = 2$ .

Construct

**$\Sigma'$  = Considering all equations of  $\Sigma$  in  $x_2, \dots, x_n, x_{n+1}$  and replacing  $x_1$  by  $x_2t$ .**

$\Sigma'$  is consistent because  $(b_2, \dots, b_n, b_{n+1})$  is the solution of  $\Sigma$  in B. So, by induction  $\Sigma'$  has a solution  $(a_2, a_3, \dots, a_n, a_{n+1})$  in A.

**$(a_1 = a_2s, a_2, \dots, a_n, a_{n+1})$  is the solution of  $\Sigma$  in A.**

## Case IV:

Suppose  $\Sigma$  has equations of the form :

$$x_s = a \text{ for } i=1,2,3,\dots,n$$

We prove this case for  $i = 1, 2$  and by the similar argument it would be true for  $i = 1, 2, \dots, n+1$ .

Rename  $x_1$  by  $x$  and  $x_2$  by  $y$  and consider

$$\Sigma = \{ x\alpha = a, xs' = xt', xs = yt, yu = b \}$$

### Subcase I:

If  $s=1$ , then  $\Sigma$  has a solution in  $A$ .

### Subcase II:

If  $s = 0$ , then construct

$$\Sigma' = \{ x\alpha = a, xs' = xt', x0 = a.0 \}$$

$$\Sigma^* = \{ yu = b, yt = a.0 \}$$

$\Sigma'$  and  $\Sigma^*$  both are consistent having solution  $b_1$  and  $b_2$  in B. So, by induction,  $\Sigma'$  and  $\Sigma^*$  having solution  $a_1$  and  $a_2$  in A, respectively.

$(a_1, a_2)$  is the solution for  $\Sigma$ .

### Subcase III:

If  $s = \alpha$ , then construct

$$\Sigma' = \{ x\alpha = a, xs' = xt' \}$$

$$\Sigma^* = \{ yu = b, yt = a \}$$

$\Sigma'$  and  $\Sigma^*$  both are consistent having solution  $b_1$  and  $b_2$  in B. So, by induction,  $\Sigma$  and  $\Sigma^*$  having solution  $a_1$  and  $a_2$  in A, respectively.

$(a_1, a_2)$  is the solution for  $\Sigma$ .

#### Subcase IV:

If  $s = \beta$

$$\Sigma = \{ x\alpha = a, xs' = xt', x\beta = yt, yu = b \}.$$

Consider

$$\Sigma' = \{ x\alpha = a, xs' = xt' \}.$$

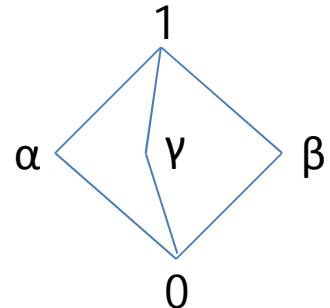
Then corresponds to every  $\sigma = \langle (s', t') \rangle$ ,  $\Sigma$  has solution in  $A$ .

So, if

$$\Sigma = \{ x\alpha = a, x\beta = yt, yu = b \}.$$

Then, corresponds to every value of  $t$  and  $u$ , we can either split  $\Sigma$  into two system in one variable or  $(a, b)$  is the solution of  $\Sigma$ .

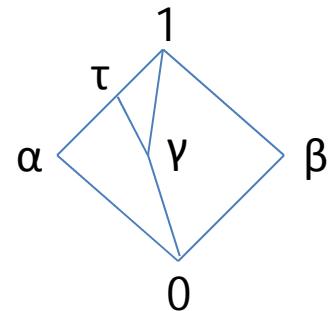
Consider the semilattice  $S = \{0, \alpha, \beta, \gamma, 1\}$ .



- ❖  $S$  is an  $S$ -act over itself.
- ❖  $S$  is not completely right pure because the principle right ideals  $\{0, \alpha\}$  and  $\{0, \beta\}$  are not linearly ordered with respect to inclusion.
- ❖  $S$  is not an almost pure  $S$ -act. Consider the system of equation

$$\Sigma = \{ x\alpha = \alpha, x\beta = 0, x\gamma = \gamma \}$$

$\Sigma$  is consistent because it has solution in extension  $T = S \cup \{\tau\}$



But  $\Sigma$  has no solution in  $S$ .

For  $S = \{0, \alpha, \beta, \gamma, 1\}$ , every almost pure  $S$ -act is absolutely pure.

For the following semilattices, every almost pure s-act is absolutely pure.

