

Algebraic manipulation detection and systems of sets in groups

Sophie Huczynska
University of St Andrews

Work described is joint with various combinations from:
Maura Paterson, Gary Mullen, Jim Davis, Chris Jefferson, Silvia
Nepřinská and others...

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Information security is concerned with the **safe and private** transmission and storage of data.

Motivating questions include:

- How can a message be sent so that we can **detect** whether it has been changed during transmission?
- If we detect that a change has occurred, can we **recover** the original message - and if so, how?
- How can we **encrypt** messages/data so that they cannot feasibly be decrypted by anyone other than the intended recipient?
- ... and many more.

Manipulation detection

In this talk, we consider an **encoding system** and how to design it to **minimise the chances** that an undetected change can occur.

Applies to various situations:

- message transmission which is subject to attack
- storage device which is subject to tampering

We will be thinking in terms of the message-sending scenario.

It is helpful to model the situation as a “game” between an **encoder** and an **adversary** who is trying to “cheat” the encoder.

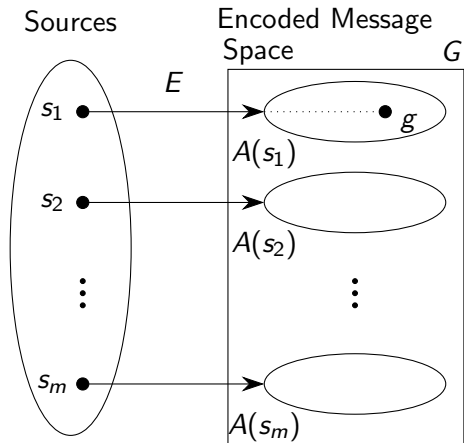
Our focus is on **algebraic manipulation detection (AMD) codes**.

We will have:

- Set S of plaintext **sources** (the messages)
- Encoded **message space** G (finite group, written additively)
- **Encoding function** E (possibly randomized) maps source $s \in S$ to some $g \in G$
- For each source $s \in S$, subset $A(s)$ of G is the set of valid encodings of s
- **Unique decodability**: $A(s) \cap A(s') = \emptyset$ if $s \neq s'$,
i.e. the sets of encodings do not overlap

Traditionally G abelian **but** the set-up is valid for non-abelian G .

Diagram



The “game”

AMD code

Adversary: chooses a value $\delta \in G \setminus \{0\}$ (their “manipulation”)

Encoder: chooses source $s \in S$

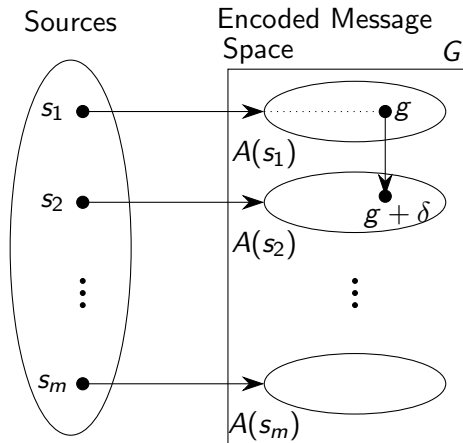
Encoder: s encoded by E to some $g \in A(s)$

Adversary: g is replaced by $g' = g + \delta$

Adversary **wins** if $g' \in A(s')$ for some $s' \neq s$

The adversary wins if they succeed in **shifting** the group element g into an element $g + \delta$ that's an encoding of a **different source**

Diagram



If message s_1 is sent and encoded as g , it will be incorrectly decoded to s_2 after this manipulation. In this case, adversary wins!

Model as system of sets

The AMD “game” can be modelled as a set-up in combinatorics.
We model the sender’s choice of message probabilistically.

Let $\{A_1, \dots, A_m\}$ be a disjoint collection of sets in G .

- Adversary chooses $\delta \in G \setminus \{0\}$
- Pick a set A_i uniformly at random (source)
- Then pick an element $d_i \in A_i$ uniformly at random (encoding)
- Adversary “wins” if $d_i + \delta \in A_j$ for some $j \neq i$

Adversary wins if δ occurs as a difference between our element in A_i and some element in A_j .

Need to look at the differences between elements of A_i and A_j .

Internal and external differences

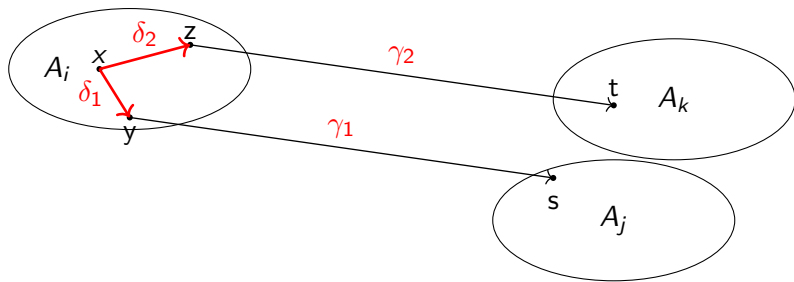
Suppose we have a disjoint family of subsets A_1, \dots, A_m of G

- For a fixed i , the differences between the elements of A_i are called **internal differences**:

$$I(A_i) := \{x - y : x, y \in A_i, x \neq y\}$$

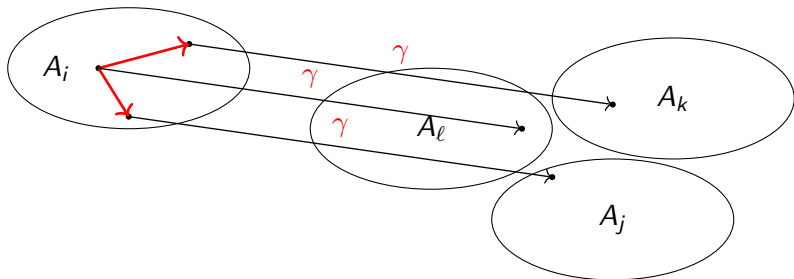
- For $i \neq j$, the differences between the elements of A_i and A_j are called **external differences**:

$$E(A_i, A_j) := \{x - y : x \in A_i, y \in A_j\}$$



In this diagram,

- δ_1 and δ_2 are **internal differences** in A_i
($x - y = \delta_1$, $x - z = \delta_2$)
- γ_1 and γ_2 are **external differences** out of A_i (to A_j , A_k resp.)
($y - s = \gamma_1$, $z - t = \gamma_2$)



For a disjoint family of sets A_1, \dots, A_m , define the **number of times** a non-zero element γ occurs as an **external difference out of A_i** by

$$N_i(\gamma) := |\{(x, y) : x - y = \gamma, x \in A_i, y \in A_j, j \neq i\}|$$

In the example above, we show all occurrences of γ as an external difference out of A_i , so $N_i(\gamma) = 3$ here.

Returning to our AMD code:

The **probability that an adversary succeeds** when they pick δ is

$$e_\delta = \frac{1}{m} \left(\frac{N_1(\delta)}{|A_1|} + \dots + \frac{N_m(\delta)}{|A_m|} \right) \quad (1)$$

- **Source i** picked with **probability $\frac{1}{m}$**
- $N_i(\delta)$ of the **possible $|A_i|$ encodings** will lead to success for an adversary who picks δ

Which codes are optimal?

We seek AMD codes that are **optimal** (from sender's perspective).

We are considering an adversary who chooses δ uniformly at random (R -strategy)

Definition

*An AMD code is (R)-optimal precisely when the **maximum success probability** of the adversary over all $\delta \in G^*$ is equal to their **average success probability**.*

Result

An AMD code is R -optimal $\Leftrightarrow e_\delta$ is constant for all $\delta \in G^$.*

We [HP] named the set systems corresponding to optimal AMD codes **reciprocally-weighted external difference families (RWEDFs)**.

Definition

An $(n, m; k_1, \dots, k_m; \ell)$ -RWEDF is a collection of disjoint subsets A_1, \dots, A_m of a group G of order n , where $|A_i| = k_i$ for all $i \in \{1, \dots, m\}$, with the property that:

$$\frac{1}{k_1} N_1(\delta) + \dots + \frac{1}{k_m} N_m(\delta) = \ell$$

for all non-zero $\delta \in G$.

In the special case when sets A_i are of equal size, this becomes

$$N_1(\delta) + \dots + N_m(\delta) = \text{constant}$$

Have been studied (abelian): **external difference families (EDFs)**.

- Let $G = (\mathbb{Z}_{10}, +)$; take $A_1 = \{4, 7, 9\}$ and $A_2 = \{0, 2, 5\}$
- Differences from A_1 to A_2 are
 $\{4 - 0 = 4, 4 - 2 = 2, 4 - 5 = -1 = 9,$
 $7 - 0 = 7, 7 - 2 = 5, 7 - 5 = 2,$
 $9 - 0 = 9, 9 - 2 = 7, 9 - 5 = 4\}$, ie $\{2, 2, 4, 4, 5, 7, 7, 9, 9\}$.
- Differences from A_2 to A_1 are their negatives, i.e.
 $\{1, 1, 3, 3, 5, 6, 6, 8, 8\}$.
- **Union of all external differences**=each nonzero element twice!
- For $\delta = 1$, the adversary's success probability is

$$\frac{1}{2} \left(\frac{N_1(\delta)}{|A_1|} + \frac{N_2(\delta)}{|A_2|} \right) = \frac{1}{2} \left(\frac{0}{3} + \frac{2}{3} \right) = \frac{1}{3}$$

- **Same** for any choice of $\delta \neq 0$.

Construction

Let G be the additive group of $GF(q)$, the finite field of order q , where q is a prime power congruent to 1 mod 4.

Let $A_1 = \{\text{the set of non-zero squares in } GF(q)\}$.

Let $A_2 = \{\text{the set of non-squares in } GF(q)\}$.

Then $\{A_1, A_2\}$ form a $(q, 2; \frac{q-1}{2}, \frac{q-1}{2}; 1)$ -RWEDF (indeed an EDF).

Special case of **cyclotomic** method - using subgroups of the multiplicative group of a finite field to make EDFs in its additive group.

What about RWEDFs which are not EDFs?

Now we would like to construct examples of RWEDFs which have different set-sizes (i.e. are not EDFs).

Example

Let $G = \mathbb{Z}_{k_1 k_2 + 1}$. The sets

$$A_1 = \{0, 1, \dots, k_1 - 1\} \text{ and } A_2 = \{k_1, 2k_1, \dots, k_1 k_2\}$$

form a $(k_1 k_2 + 1, 2; k_1, k_2; \frac{1}{k_1} + \frac{1}{k_2})$ -RWEDF.

Can prove: this gives an AMD code whose success probability is **as small as possible** when $m = 2$ for the given group size.

Definition

A difference set in a group G is a set $D \subseteq G$ such that, when we take all pairwise *internal differences* between the elements of D , every non-identity group element occurs a *fixed number* λ of times.

Result

Let G be a group of order n , and let $\mathcal{A} = \{A_1, A_2\}$ partition G . Then \mathcal{A} is an RWEDF $\Leftrightarrow A_1$ and A_2 are difference sets.

Example: Let $G = \mathbb{Z}_7$. Let $A_1 = \{1, 2, 4\}$ and $A_2 = \{0, 3, 5, 6\}$. Then $\{A_1, A_2\}$ is a $(7, 2; 3, 4; \frac{7}{6})$ -RWEDF. For any δ , adversary's success probability is $\frac{7}{12}$.

Motivating questions for RWEDFs

Observe that all the examples we have seen so far have 2 sets.

Question

Can we get examples with *more than 2 sets*, ie $m > 2$?

The constant ℓ in the definition is in \mathbb{Q} but not necessarily in \mathbb{Z} .

Question

Can we obtain constructions for RWEDFs with *integer ℓ* ?

One way to guarantee integer ℓ would be if $k_i \mid N_i$ ($1 \leq i \leq m$).

We must have $N_i(\delta) \leq k_i$ for $\delta \in G \setminus \{0\}$ - so this would mean $N_i(\delta) = 0$ or k_i .

Defining a New Property

Motivated by this condition for RWEDFs with integer ℓ , we define the following general property.

Let G be a finite group and let \mathcal{A} be a collection $\{A_1, A_2, \dots, A_m\}$ of disjoint subsets of G with sizes k_1, k_2, \dots, k_m respectively.

Definition

We shall say \mathcal{A} has the *bimodal property* if for all $\delta \in G^*$ we have $N_j(\delta) \in \{0, k_j\}$ for $j = 1, 2, \dots, m$.

In other words: for each $\delta \in G^*$, either δ **never** occurs as a difference between A_i and some other A_j , or else for **every** $a_i \in A_i$ there is an $a_j \in A_j$ ($i \neq j$) s.t. $\delta = a_i - a_j$.

Not every collection of bimodal sets will be an RWEDF...

Example

Let $G = \mathbb{Z}_{10}$ and take $A_1 = \{1, 6\}$, $A_2 = \{3, 8\}$ and $A_3 = \{4, 9\}$.
Then $A = \{A_1, A_2, A_3\}$ is bimodal but not an RWEDF.

For $i = 3$ we have $N_3(1) = N_3(3) = N_3(6) = N_3(8) = 2 = k_3$ while

$$N_3(2) = N_3(4) = N_3(5) = N_3(7) = N_3(9) = 0.$$

Similar calculations for $N_1(\delta)$ and $N_2(\delta)$ confirm A has the bimodal property **but** 5 never occurs as an external difference.

...and not every RWEDF with integer ℓ will be bimodal - though some always will:

Result

An $(n, m; k_1, \dots, k_m; \ell)$ -RWEDF with $\ell \in \mathbb{Z}$ and $\{k_1, \dots, k_m\}$ pairwise coprime is bimodal.

Understanding bimodal sets

This opens up two quite **distinct questions**:

- Can families of sets with the bimodal property in finite (abelian) groups be algebraically characterized?
- Can we find bimodal families of sets which are RWEDFs?

Result

Let H be a subgroup of an abelian group G . If $\mathcal{C} = \{C_1, \dots, C_m\}$ is a collection of cosets of H , the \mathcal{C} has the bimodal property.

Proof: For fixed i and $1 \leq j \leq m$ with $i \neq j$, the sets $C_i - C_j$ comprise $m - 1$ distinct cosets of H . For any $\delta \in C_i - C_j$ and every $x \in C_i$, \exists a unique $y \in C_j$ s.t. $x - y = \delta$. However, for any $\delta \in G \setminus \cup_{j \neq i} (C_i - C_j)$, δ occurs 0 times as a difference out of C_i .

Cosets are such a “natural” example that you may guess they are the only non-trivial collection of sets with the bimodal property, but in fact a much richer landscape emerges.

Definition

Let A be a subset of a finite abelian group G . We define the *internal difference group* H of A to be the subgroup of G generated by all $x - y$ where $x, y \in A$, ie $H = \langle I(A) \rangle$.

The group H has the property that A is contained in a single coset of H , and is the smallest subgroup of G with this property.

Bimodal property and cosets

For disjoint subsets $\{A_1, \dots, A_m\}$ of our group G , we will let:

- $A = \cup_{i=1}^m A_i$
- $B_i = A \setminus A_i$ for any $1 \leq i \leq m$
- $H_i = \langle I(A_i) \rangle$

Result

Let G be a finite abelian group and let $\mathcal{A} = \{A_1, \dots, A_m\}$ be a collection of disjoint subsets of G . Then \mathcal{A} has the bimodal property if and only if for each i the set B_i is a union of cosets of the subgroup H_i .

Definition

If a finite group G has subgroups S_1, \dots, S_m with the property that $S_1 \setminus \{0\}, \dots, S_m \setminus \{0\}$ partition $G \setminus \{0\}$, then the collection of subgroups is called **a partition of G** .

Example

Let $G = \mathbb{Z}_3 \times \mathbb{Z}_3$.

A partition of G is given by:

$$S_1 = \{(0,0), (1,1), (2,2)\}, S_2 = \{(0,0), (0,1), (0,2)\},$$

$$S_3 = \{(0,0), (1,2), (2,1)\}, S_4 = \{(0,0), (1,0), (2,0)\}.$$

Subgroup partitions give bimodal sets

Let S_i^* denote $S_i \setminus \{0\}$.

Result

If a finite abelian group G has subgroups S_1, \dots, S_m forming a partition of G , then $\{S_1^, \dots, S_m^*\}$ has the bimodal property.*

Proof: For each i , the internal difference group of S_i^* is S_i itself. So the union $\cup_{j \neq i} S_j^*$ is $G \setminus S_i$, a union of cosets of S_i .

Have seen: **cosets** and **subgroup partitions** - what is the general landscape for collections of bimodal sets?

Impressionistic idea of the situation

- Collections of bimodal sets in finite abelian groups are in some sense a “blend” of the coset and subgroup partition examples we’ve seen.
- Let $r_{\mathcal{A}}$ be the number of A_i with $|A_i| < |H_i|$.
Key structural differences when $r_{\mathcal{A}} \geq 2$, $= 1$ and $= 0$.
- When $r_{\mathcal{A}} \geq 2$, the sets $A_1, \dots, A_{r_{\mathcal{A}}}$ occur together in a very tightly-structured way: like an “inflated” group partition.
- This imposes considerable structure on the remaining members of \mathcal{A} (coset part).
- The cases with $r_{\mathcal{A}} = 1$ and $= 0$ are comparable but have a simpler description.

Some technical points for our main structure result

- Recall $|A_i| \leq |H_i|$ for all $1 \leq i \leq m$.
Wlog, we label the sets such that
 - $|A_i| < |H_i|$ for $i = 1, \dots, r_{\mathcal{A}}$
 - $|A_i| = |H_i|$ for $i = r_{\mathcal{A}} + 1, \dots, m$.
- Following literature, a collection of sets F_1, \dots, F_k with the property that $F_i \cap F_j = D$ for all for all $i \neq j$ is said to be a ***k*-star with kernel D** .
- Helpful to shift to **“canonical position”**: a translation guaranteeing that instead of cosets of certain subgroups, we are dealing with the subgroups themselves.

Structural result

Let $\mathcal{A} = \{A_1, \dots, A_m\}$ be a bimodal collection of disjoint subsets of an abelian group G with $r_{\mathcal{A}} \geq 2$, in canonical position.

Result

- *The internal difference groups $H_1, \dots, H_{r_{\mathcal{A}}}$ form an $r_{\mathcal{A}}$ -star with kernel $D_{\mathcal{A}}$ (a subgroup of G), and for each i with $1 \leq i \leq r_{\mathcal{A}}$ we have $A_i = H_i \setminus D_{\mathcal{A}}$.*
- *Any set A_i with $i > r_{\mathcal{A}}$ is a coset of a subgroup of $D_{\mathcal{A}}$*
- *If H denotes the group $H_1 + H_2 + \dots + H_{r_{\mathcal{A}}}$, then $H \setminus D_{\mathcal{A}}$ is contained in A . Furthermore, the sets in \mathcal{A} can be labelled such that for some k with $r_{\mathcal{A}} \leq k \leq m$ we have that $H \setminus D_{\mathcal{A}}$ is partitioned by A_1, \dots, A_k .*
- *If $k < m$ then the sets A_i with $i > k$ arise from a subdivision of cosets of H .*

We also have the other way round...

Let G be an abelian group, and for $t \geq 2$ let H_1, \dots, H_t be distinct subgroups of G forming a t -star with kernel D , such that

$|H_i : D| > 2$ for i with $1 \leq i \leq t$.

Let $H = H_1 + \dots + H_t$.

Result

Let \mathcal{A} consist of the following subsets of G :

- all subsets of the form $A_i = H_i \setminus D$ for i with $1 \leq i \leq t$;
- all cosets of D that are subsets of H , but are not in $\cup_{i=1}^t H_i$;
- for any number of cosets of H , all the cosets of D that lie within those cosets of H .

Then \mathcal{A} is a bimodal collection of subsets of G with $r_{\mathcal{A}} = t$ in canonical position.

Returning to RWEDFs, we can prove the following for **any** (not necessarily abelian) finite group:

Result

If a finite group G of order n has subgroups S_1, \dots, S_m forming a partition of G , then $\{S_1^, \dots, S_m^*\}$ is a (bimodal) RWEDF.*

- This gives a wealth of new RWEDF/EDF examples, in both abelian and non-abelian groups.

From the literature, groups which admit a subgroup partition include:

- elementary abelian p -groups of order $\geq p^2$, for p prime
- Frobenius groups (eg dihedral group D_{2n} with n odd)
- groups of Hughes-Thompson type
- groups isomorphic to $PGL(2, p^h)$ with p an odd prime

Example

- Let $G = \mathbb{Z}_3 \times \mathbb{Z}_3$.
- We use the subgroup partition from the earlier slide, removing the zero element.
- Let $A_1 = \{(1, 1), (2, 2)\}$, $A_2 = \{(0, 1), (0, 2)\}$,
 $A_3 = \{(1, 2), (2, 1)\}$ and $A_4 = \{(1, 0), (2, 0)\}$.
- For non-zero $\delta \in G$, $N_i(\delta) = 2$ for $\delta \notin A_i$ and $N_i(\delta) = 0$ for $\delta \in A_i$ (for each $1 \leq i \leq 4$).
- \mathcal{A} forms a $(9, 4; 2, 2, 2, 2; 3)$ -RWEDF (indeed, this is an EDF).

This is an example of a more general construction we have in elementary abelian p -groups using **vector space partitions**.

Example

Let n be odd, and let D_{2n} be the *dihedral group* given by the presentation

$$\langle r, s \mid r^n = 1, s^2 = 1, rs = sr^{-1} \rangle$$

A partition is given by $S_i = \langle sr^{i-1} \rangle$ for $1 \leq i \leq n$ and $S_{n+1} = \langle r \rangle$.
Here $|S_1| = \dots = |S_n| = 2$ and $|S_{n+1}| = n$.
For D_{10} this yields a $(10, 6; 1, 1, 1, 1, 1, 4; 5)$ -RWEDF.

Nonabelian RWEDF partition example

Example

Let G be the *Heisenberg group* modulo 3, ie the group of 3×3 upper triangle matrices with entries from $GF(3)$ that have 1s on the main diagonal.

Each element of G has the form $\begin{pmatrix} 1 & a & b \\ 0 & 1 & c \\ 0 & 0 & 1 \end{pmatrix}$ and each

non-identity element has order 3.

$|G| = 27$ and its order 3 subgroups partition its non-identity elements.

This will give an EDF with 13 sets of size 2.

A stronger security model

Question: is it always realistic to assume that an adversary will not know which message (source) is being sent?

We may wish to consider a **stronger** security model, in which the adversary knows the source **before** they choose their δ .

Strong AMD code

Encoder: chooses a source $s \in S$

Adversary: is given source s

Adversary: chooses some $\delta \in G \setminus \{0\}$

Encoder: source is encoded by E to $g \in A(s)$

Adversary: g is replaced by $g' = g + \delta$

Adversary **wins** if $g' \in A(s')$ for some $s' \neq s$.

In a **strong** AMD code, the adversary **learns** s before choosing δ .

Strong AMD codes and set systems

What mathematical structures correspond to optimal strong AMD codes?

At present: **strong EDFs** are used.

These require a condition **for each possible i** :

Definition

A **strong external difference family** in an abelian group G of order n is a collection of **disjoint sets** A_1, \dots, A_m of G , each of size k , such that when we take all **external differences** from **any A_i** to $\cup_{j \neq i} A_j$, every non-identity group element occurs a fixed number ℓ of times.

We write this as an (n, m, k, ℓ) -SEDF.

In fact, we have seen **examples of SEDFs** earlier in the talk:

- In $G = \mathbb{Z}_5$, the sets $\{1, 4\}$ and $\{2, 3\}$ form an SEDF.
- Let G be the additive group of $GF(q)$ where q is a prime power congruent to 1 mod 4; the set of non-zero squares and the set of non-squares form an SEDF.
- In $G = \mathbb{Z}_{k^2+1}$, the sets

$$A_1 = \{0, 1, \dots, k-1\} \text{ and } A_2 = \{k, 2k, \dots, k^2\}$$

form a $(k^2 + 1, 2, k, 1)$ -SEDF.

The SEDF landscape: existence

SEDFs are known to **exist** for the following (n, m, k, ℓ) :

- (a) $(k^2 + 1, 2, k, 1)$: $G = \mathbb{Z}_{k^2+1}$, Paterson/Stinson
- (b) $(v, 2, \frac{v-1}{2}, \frac{v-1}{4})$ where $v \equiv 1 \pmod{4}$ and an appropriate partial difference set exists: Davis/Huczynska/Mullen and Huczynska/Paterson
- (c) $(q, 2, \frac{q-1}{4}, \frac{q-1}{16})$ where $q = 16t^2 + 1$ is a prime power and $G = GF(q)$: Bao/Wei/Zhang
- (d) $(q, 2, \frac{q-1}{6}, \frac{q-1}{36})$ where $q = 108t^2 + 1$ is a prime power and $G = GF(q)$: Bao/Wei/Zhang

First SEDF with $m > 2$

Until the start of 2018, **no SEDFs** were known with $m \neq 2$.
Then the first with $m > 2$ was found - **independently** by two sets of authors.

(243, 11, 22, 20)-SEDF in \mathbb{Z}_3^5

- Cyclotomic construction [Wen, Yang, Feng]
- Action of M_{11} on $\text{PG}(4, 3)$ [Jedwab, Li]

This is still the **only known** SEDF with more than 2 sets!

Non-abelian SEDFs

One theme of the bimodal work was the emergence of **non-abelian** RWEDF examples.

Recently, we [HJN] obtained the first construction for a family of **non-abelian SEDFs**:

Theorem

Let $k > 1$ be odd. In D_{k^2+1} , the dihedral group of order $n = k^2 + 1$, there exists a $(k^2 + 1, 2, k, 1)$ -SEDF in G . Specifically, in

$$\langle r, s \mid r^{n/2} = 1, s^2 = 1, rs = sr^{-1} \rangle$$

we can take $\{A_1, A_2\}$ where

- $A_1 = \{r^i : 0 \leq i \leq \frac{k-1}{2}\} \cup \{sr^j : 0 \leq j \leq \frac{k-3}{2}\}$.
- $A_2 = \{r^{ik} : 1 \leq i \leq \frac{k-1}{2}\} \cup \{sr^{\frac{k(2j+1)-1}{2}} : 0 \leq j \leq \frac{k-1}{2}\}$.

There are many avenues to explore further in this area.

- Beyond group partitions, which collections of bimodal sets **guarantee RWEDFs**?
- Obtain a **combinatorial characterization** of RWEDFs with integer ℓ .
- Constructions for RWEDFs with integer ℓ which are **not bimodal**?
- **Fine-tune** our constructions to yield smallest possible success probabilities.
- The **strong** model for RWEDFs.
- Further constructions in **nonabelian** groups.
- Specific connections with **Frobenius** groups?

Thank you for listening!