

# ***S-acts, coherency and stability***

York Semigroup, Spring 2013

**Victoria Gould**  
**University of York**

## Monoids and acts

A **(right)  $S$ -act** is a set  $A$  together with a map

$$A \times S \rightarrow A, (a, s) \mapsto as$$

such that for all  $a \in A, s, t \in S$

$$a1 = a \text{ and } (as)t = a(st).$$

Right ideals (including  $S$ ) are  $S$ -acts.

Let  $A$  be an  $S$ -act.

For any  $s \in S$ , we have a unary operation  $\rho_s$  on  $A$  given by  $a \mapsto as$  and a morphism  $\phi : S \rightarrow \mathcal{T}_A$  given by  $s \mapsto \rho_s$ .

Conversely, if  $\phi : S \rightarrow \mathcal{T}_B$  is a morphism for a set  $B$ , then  $B$  is an  $S$ -act with  $bs = b(s\phi)$ .

Consequently,  $S$ -acts are **representations of monoids by mappings of sets**.

## Elementary observations for $S$ -acts

- An  **$S$ -morphism** from  $A$  to  $B$  is a map  $\alpha : A \rightarrow B$  with  $(as)\alpha = (a\alpha)s$  for all  $a \in A, s \in S$ .
- $S$ -acts and  $S$ -morphisms form a category - products are products, coproducts are disjoint unions
- We have usual definitions of **free**, **projective**, **injective**, etc. including variations on **flat**.
- Free  $S$ -acts are disjoint unions of copies of  $S$ .

## Elementary observations for $S$ -acts cont.

- A **congruence**  $\rho$  on  $A$  is an equivalence relation such that

$$a \rho b \Rightarrow as \rho bs$$

for all  $a, b \in A$  and  $s \in S$ .

- $\rho$  is **finitely generated** if  $\rho$  is the smallest congruence containing a finite set  $H \subseteq A \times A$ .
- If  $\rho$  is a congruence on  $A$  then  $A/\rho$  is an  $S$ -act; all monogenic  $S$ -acts are of the form  $S/\rho$ .
- An  $S$ -act  $A$  is **finitely generated** if

$$A = a_1S \cup \dots \cup a_nS$$

for some  $a_i \in A$  and **finitely presented** if

$$A \cong F_n/\rho$$

for some finitely generated free  $F_n$  and finitely generated congruence  $\rho$ .

# First order languages and $L_S$

A (first order) language  $L$  has alphabet:

variables, connectives (e.g.  $\neg, \vee, \wedge, \rightarrow$  etc.), quantifiers  $(\forall, \exists)$ ,  $=$ ,  
brackets, commas  
and some/all of symbols for constants, functions and relations.

There are rules for forming well formed formulae (wff); a **sentence** is a wff  
with no free variables (i.e. all variables are governed by quantifiers).

**The language  $L_S$**  has:

no constant or relational symbols (other than  $=$ )  
for each  $s \in S$ , a unary function symbol  $\rho_s$ .

**A point of convenience** Let us agree to abbreviate  $x\rho_s$  in wffs of  $L_S$  by  $xs$ .

# $L_S$ -structures and $S$ -acts

## Examples

$\neg(xs = xt)$  is a wff but not a sentence,

$(\forall x)(\neg(xs = xt))$  is a sentence,

$(\exists \vee xsx$  is not a wff.

An  **$L$ -structure** is a set  $D$  equipped with enough distinguished elements (constants), functions and relations to ‘interpret’ the abstract symbols of  $L$ .

An  $L_S$ -structure is simply a set with a unary operation for each  $s \in S$ .

Clearly an  $S$ -act  $A$  is an  $L_S$ -structure where we interpret  $\rho_s$  by the map  $x \mapsto xs$ .

# Model theory

A **theory** is a set of sentences in a first order language.

**Model theory** provides a range of techniques to study algebraic and relational structures etc. via properties of their associated **languages** and **theories**.

**Model theory of  $R$ -modules** is a well developed subject area.

**Model theory of  $S$ -acts** - much less is known - authors include **Ivanov, Mustafin, Stepanova** .

**Stability** is an area within model theory, introduced by **Morley 62**. Much of the development of the subject is due to **Shelah**; the definitive reference is **Shelah 90** though (quote from Wiki) *it is notoriously hard even for experts to read*.

## Model theory lite: axiomatisability

**Definition** A class  $\mathcal{A}$  of  $L_S$ -structures is **axiomatisable** if there is a theory  $\Sigma$  such that for any  $L_S$ -structure  $A$ , we have  $A \in \mathcal{A}$  if and only if every sentence of  $\Sigma$  is true in  $A$ , i.e.  $A$  is a *model* of  $\Sigma$ .

**Example** Let  $\Sigma_S$  be the theory

$$\Sigma_S = \{(\forall x)((xs)t = x(st)) : s, t \in S\} \cup \{(\forall x)(x1 = x)\}.$$

Then  $\Sigma_S$  axiomatises the class of  $S$ -acts (within all  $L_S$ -structures).

**Example** Let  $\Pi_S$  be the theory

$$\Sigma_S \cup \{(\exists x)(\neg(xs = xt)) : s, t \in S, s \neq t\}.$$

Then  $\Pi_S$  axiomatises the faithful  $S$ -acts.

## Model theory of $S$ -acts

**Mustafin 88** There exists a monoid  $S$  and an  $S$ -act  $A$  such that  $\text{Th}(A)$  is not stable. *This contrasts with the situation for modules.*

**Our aim today** To look at finitary properties for monoids arising from existence and stability of the model companion of  $\Sigma_S$ .

A **finitary property** for a monoid is one held by finite monoids, e.g.  $S$  is **weakly right noetherian** if  $S$  has the ascending chain condition on right ideals.

Some classes of  $S$ -acts (such as the projectives and injectives) are axiomatisable iff  $S$  satisfies some finitary conditions.

## Existentially closed $S$ -acts

Let  $A$  be an  $S$ -act. An **equation** over  $A$  has the form

$$xs = xt, \quad xs = yt \text{ or } xs = a$$

where  $x, y$  are variables,  $s, t \in S$  and  $a \in A$ . **Inequations** look like

$$xs \neq xt, \quad xs \neq yt \text{ or } xs \neq a.$$

A set of equations and inequations is **consistent** if it has a solution in some  $S$ -act  $B \supseteq A$ .

**Definition**  $A$  is **existentially closed** if every finite consistent set of equations and inequations over  $A$  has a solution in  $A$ .

Let  $\mathcal{E}$  denote the class of existentially closed  $S$ -acts.

**Question** When is  $\mathcal{E}$  axiomatisable?

## Model companions

**Definition** Let  $T, T^*$  be theories in a first order language  $L$ . Then  $T^*$  is a **model companion** of  $T$  if every model of  $T$  embeds into a model of  $T^*$  and vice versa, and embeddings between models of  $T^*$  are elementary embeddings.

**Theorem Wheeler 76**  $\Sigma_S$  has a model companion  $\Sigma_S^*$  precisely when  $\mathcal{E}$  (the class of existentially closed  $S$ -acts) is axiomatisable and in this case,  $\Sigma_S^*$  axiomatises  $\mathcal{E}$ .

**Question** When does  $\Sigma_S^*$  exist? i.e. When is  $\mathcal{E}$  axiomatisable?

## Right coherent monoids

**Definition**  $S$  is **right coherent** if every finitely generated  $S$ -subact of every finitely presented  $S$ -act is finitely presented.

Let  $A$  be an  $S$ -act and let  $z \in A$ . Put

$$\mathbf{r}(z) = \{(u, v) \in S \times S : zu = zv\}$$

and notice that  $\mathbf{r}(z)$  is a right congruence on  $S$ .

**Theorem Wheeler 76, G 87, 92, Ivanov 92** The f.a.e. for  $S$ :

- ①  $\Sigma_S^*$  exists;
- ②  $S$  is right coherent;
- ③ every finitely generated  $S$ -subact of every  $S/\rho$ , where  $\rho$  is finitely generated, is finitely presented;
- ④ for every finitely generated right congruence  $\rho$  on  $S$  and every  $a, b \in S$  we have  $\mathbf{r}([a])$  is finitely generated, and  $[a]S \cap [b]S$  is finitely generated.

## Right coherent and right noetherian monoids

**Definition**  $S$  is **right noetherian** if every right congruence is finitely generated.

**Fact** If  $S$  is right noetherian, it is weakly right noetherian.

**Theorem Normak 77** If  $S$  is right noetherian, it is right coherent.

**Example Fountain 92** There exists a weakly right noetherian  $S$  which is not right coherent.

## Right coherent and right noetherian monoids

Definitions  $S$  is

- ① **regular** if  $\forall a \in S \exists x \in S, a = axa$ ;
- ② **inverse** if  $S$  is regular and  $ef = fe$  for all  $e = e^2, f = f^2 \in S$ ;
- ③ **Clifford** if  $S$  is regular and  $ae = ea$  for all  $a, e = e^2 \in S$ .

## Right coherent and right noetherian monoids: Examples

Results variously due to Gould, Hartmann, Ruskuc, Yang

- ① **1992** The free commutative monoid  $\mathcal{FC}(X)$  on  $X$ ; *this follows from Rédei's theorem that if  $X$  is finite, then  $\mathcal{FC}(X)$  is noetherian, so coherent from Normak; an easy argument then gives the infinite case*
- ② **1992** Clifford monoids;
- ③ **2012** the free monoid  $X^*$ ;
- ④ **2005, 2011** weakly right noetherian regular monoids so the **bicyclic monoid** and  $BR(G, \theta)$ ;
- ⑤ **2012** the *double Bruck-Reilly* with identity adjoined;
- ⑥ **2013** primitive inverse monoids,  $\mathcal{B}^0(M, I)$  where  $M$  is right coherent;
- ⑦ **2012** Free left ample
- ⑧ **2012** Free inverse monoids are NOT right coherent.

## Stability: a bit of background

Stability properties for theories - **stable, superstable and totally transcendental** - arose from the question of how many models a theory has of any given cardinality.

**Shelah 78** Showed that a non-superstable theory, has  $2^\lambda$  models of cardinality  $\lambda$  for any  $\lambda > \max\{\aleph_0, |T|\}$ .

The philosophy then is that, in these cases, there are too many models to attempt to classify (e.g. by means of a sensible structure theorem).

It is reasonable therefore for the algebraist to consider for a given (axiomatisable) class of algebras 'how stable' is the theory associated with it.

## A quick view of stability properties

We will be looking at stability properties of  $\Sigma_S^*$ .

We have ascertained that for  $\Sigma_S^*$  to exist,  $S$  must be right coherent.

Assume now that  $S$  is right coherent.

Is  $\Sigma_S^*$  **stable**? If so, when is  $\Sigma_S^*$  **superstable**?

...and when is  $\Sigma_S^*$  **totally transcendental**?

# Types

Let  $A$  be an  $S$ -act. Then  $L_S(A)$  is the language  $L_S$  augmented with symbols representing the elements of  $A$ .

$A$  is an  $S$ -subact of  $\mathbf{M}$  where  $\mathbf{M}$  is the ‘monster model’ of  $\Sigma_S^*$ . Let  $c \in \mathbf{M}$ . Then

$$\text{tp}(c/A) = \{\phi(x) \in L_S(A) : \mathbf{M} \models \phi(c)\}$$

is a **type over  $A$** .

The **Stone space**  $S(A)$  of  $A$  is the collection of all types over  $A$  equipped with the topology with basis

$$\langle \phi(x) \rangle = \{p \in S(A) : \phi(x) \in p\}.$$

## Stability conditions

**Definition** For an infinite cardinal  $\kappa$  we say  $\Sigma_S^*$  is:

- **$\kappa$ -stable** if for all  $S$ -acts  $A$  with  $|A| \leq \kappa$  we have  $|S(A)| \leq \kappa$ ;
- **stable** if  $\Sigma_S^*$  is  $\kappa$ -stable for *some*  $\kappa$ ;
- **superstable** if  $\Sigma_S^*$  is  $\kappa$ -stable for all  $\kappa \geq 2^{|T|}$  where  $T = \Sigma_S^*$ .

**Fact Lascar, 76** A complete theory is superstable if and only if every type  $p$  has U-rank  $U(p) < \infty$ .

**Definition** A complete theory is **totally transcendental** if and only if every type  $p$  has Morley rank  $M(p) < \infty$ .

**Fact** For any type  $p$  over a theory  $T$  we have  $U(p) \leq M(p) \leq \infty$ .

## Stability of $\Sigma_S^*$

Let  $\mathcal{RI}$  and  $\mathcal{RC}$  denote the lattices of right ideals and right congruences on  $S$ .

Let  $A$  be an  $S$ -act. An  $A$ -triple  $(I, \rho, f)$  is a triple such that:

$$I \in \mathcal{RI}, \rho \in \mathcal{RC}, I \text{ is } \rho\text{-saturated}$$

and

$$f : I \rightarrow A \text{ is an } S\text{-morphism with } \text{Ker } f = \rho \cap (I \times I).$$

Let  $T(A)$  denote the set of all  $A$ -triples.

**Theorem Fountain, G 87/08**  $S(A)$  is in bijective correspondence with  $T(A)$  under  $p \mapsto (I_p, \rho_p, f_p)$ .

## Stability and superstability of $\Sigma_S^*$

**Theorem Fountain, G 87/08, Ivanov 92**

- ①  $\Sigma_S^*$  is stable.
- ②  $\Sigma_S^*$  is superstable if and only if  $S$  is weakly right noetherian.

We prove this by using a notion of rank on  $\rho$ -saturated right ideals for a right congruence  $\rho$  so get a value of U-rank of a type in algebraic terms.

## Weakly right noetherian (wrn)

- ① Groups are wrn.
- ② Semilattices are wrn iff they have the ascending chain condition (as posets) and no infinite antichains (folklore).
- ③ Bicyclic monoid  $B = \mathbb{N}^0 \times \mathbb{N}^0$  with binary operation

$$(a, b)(c, d) = (a - b + t, d - c + t) \text{ where } t = \max\{b, c\}$$

is inverse and wrn.

- ④ Free inverse monoids on non trivial  $X$  are not wrn.
- ⑤  $S$  regular, wrn implies  $S$  right coherent.
- ⑥ Brandt semigroups with identity adjoined are right coherent but need not be wrn.

## Total transcendence

A theory  $T$  is **totally transcendental** if and only every type  $p$  has Morley rank  $M(p)$ . These are the best theories in terms of stability.

A totally transcendental theory is superstable.

**Fountain, G 87/08** characterised those (right coherent)  $S$  such that for every type  $p$ , we have  $U(p) = M(p) < \infty$ .

**Note** that if  $M(p) < \infty$  for a type over a theory of modules, then  $U(p) = M(p)$  - is this true for  $S$ -acts?

## The finite type topology in $\mathcal{RC}$

We define a topology on  $\mathcal{RC}$  by means of a basis.

Let  $\nu \in \mathcal{RC}$  be finitely generated and let  $K \subseteq (S \times S) \setminus \nu$  be finite.

Put

$$[\nu, K] = \{\rho \in \mathcal{RC} : \nu \subseteq \rho \subseteq (S \times S) \setminus K\}.$$

The **finite type topology** has sets  $[\nu, K]$  as a basis.

The **Cantor-Bendixon** rank of a point in topological space measures how far a point is from being isolated.

The **S-rank** of  $\rho \in \mathcal{RC}$  is the Cantor-Bendixon rank of  $\rho$  with respect to the finite type topology.

## *S*-rank

We make this explicit by defining subsets  $\mathcal{C}^\alpha$  of  $\mathcal{RC}$  for each ordinal  $\alpha$ , as follows:

- (I)  $\mathcal{C}^0 = \mathcal{C}$ ;
- (II) if  $\alpha$  is a limit ordinal, then

$$\mathcal{C}^\alpha = \bigcap \{\mathcal{C}^\beta : \beta < \alpha\};$$

- (III)  $\rho \in \mathcal{C}^{\alpha+1}$  if and only if  $\rho \in \mathcal{C}^\alpha$  and for all subsets of finite type  $[\nu, K]$  with  $\rho \in [\nu, K]$ , there exists  $\theta \in \mathcal{C}^\alpha$  with

$$\theta \in [\nu, K], \theta \neq \rho.$$

The **S-rank**  $S(\rho)$  of  $\rho \in \mathcal{RC}$  is  $\infty$  if  $\rho \in \mathcal{C}^\alpha$  for all  $\alpha$ , and otherwise  $S(\rho) = \alpha$  where  $\rho \in \mathcal{C}^\alpha \setminus \mathcal{C}^{\alpha+1}$ . If  $S(\rho) < \infty$ , then we say that  $\rho$  **has S-rank**.

*S* is **ranked** if every  $\rho \in \mathcal{RC}$  has *S*-rank.

# Ranked semigroups

**Theorem G 05**  $\Sigma_S^*$  is totally transcendental if and only if  $S$  is wrn and every  $\rho \in \mathcal{RC}$  has  $S$ -rank.

- ① If  $S$  is right noetherian, then  $S$  is ranked.
- ② If  $S$  inverse and ranked, then  $S$  is wrn.
- ③ Let  $S$  be ranked. Then any maximal subgroup and monoid  $\mathcal{J}$ -class is ranked.
- ④ A ranked monoid cannot contain a bicyclic  $\mathcal{J}$ -class.
- ⑤ The chain  $C$

$$e_0 > e_1 > \dots$$

is right coherent, weakly right noetherian but not ranked, so  $\Sigma_C^*$  is superstable but not totally transcendental.

- ⑥  $\mathbb{Z}$  under  $+$  is (right) noetherian, right coherent and ranked so  $\Sigma_{\mathbb{Z}}^*$  is t.t., but there is a type  $p$  with  $M(p) \neq U(p)$ .

## Questions

- ① For an arbitrary  $S$ , does ranked imply wrn?
- ② Connections of right coherency to products of (weakly, strongly) flat *left S-acts*?
- ③ Ranked groups?
- ④ Structure of  $S$ -acts where  $S$  satisfies suitable finitary properties.